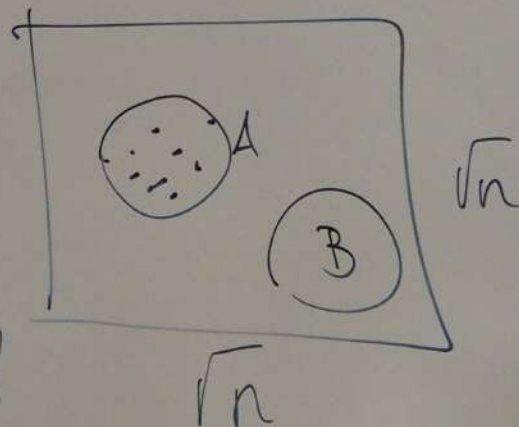


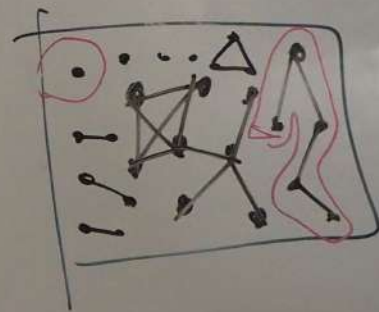
Connectivity of RGGs ($\text{Poisson}[0, \sqrt{n}]^2$)

$$G(X_n, r), \quad r = C\sqrt{\log n}$$

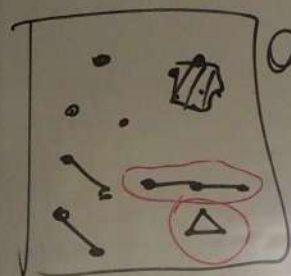


Thm: The threshold for having an isolated vertex and being connected are equal.

Lemma 1: Let $c > 0$. Whp, $G(X_n, c\sqrt{\log n})$ does not have 2 components of Euclidean diameter $\geq C\sqrt{\log n}$ for large enough $C = C(c) > 0$.



Lemma 2: let C as in lemma 1. Whp the graph $G = (V_n, \sqrt{\frac{\log n - \frac{1}{2} \log \log n}{\#}})$ does not have a component with > 1 point and Euclidean diameter $\leq C \sqrt{\log n}$.



Proof: Suppose H is such a component. Case 1: Euclidean diameter of H is $\leq \frac{(\log \log n)^2}{\sqrt{\log n}} = \eta$. let $x \in H$.

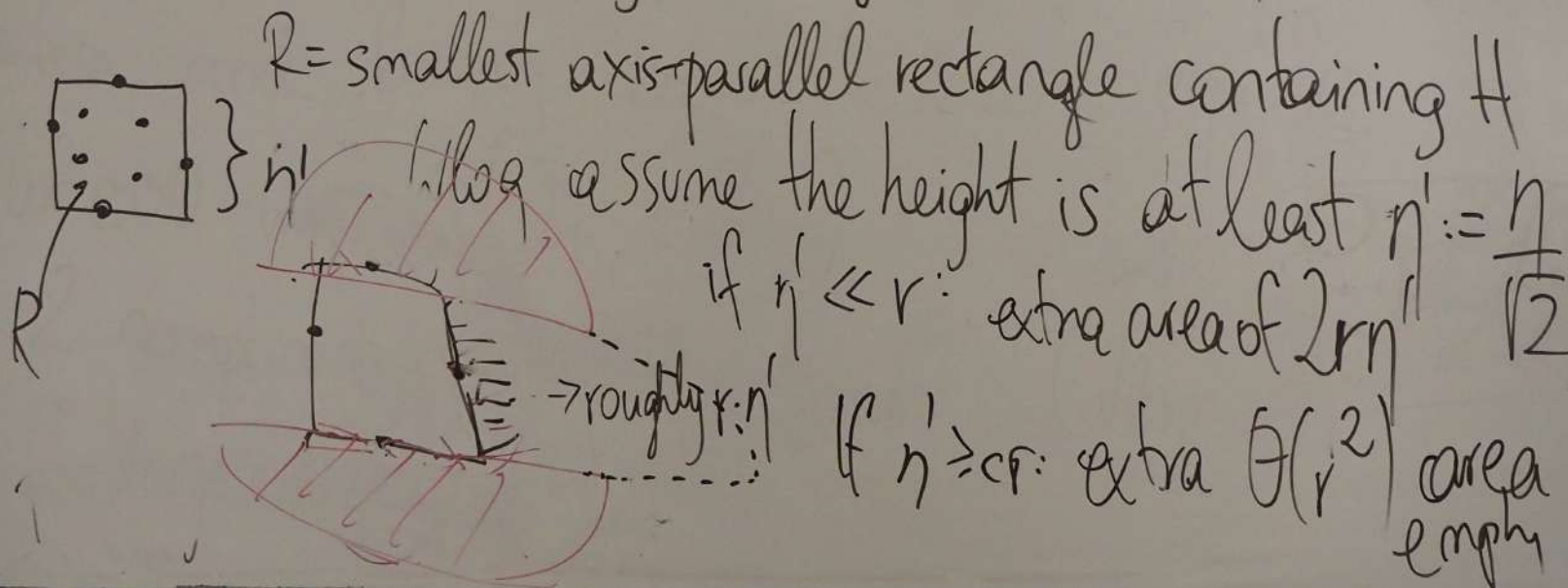


Exercise: fix x . $P(|B(x, \eta) \cap X_n| \geq 1, |B(x, r) \setminus B(x, \eta) \cap X_n| = 0) = o(1/n)$

→ Wlog, there is no such vertex x .

Case 2: Euclidean diameter of $H \geq \eta$.

Consider the following rectangle:



Hence, empty area is at least

$$r^2 \pi + 2rn'(1+o(1)) \geq r^2 \pi + \sqrt{2}rn(1+o(1))$$
$$= r^2 \pi + (\log \log n)^2$$

$$P(\text{this area empty}) = \exp\left(-\log n + \frac{1}{2} \log \log n - (\log \log n)^2\right) = O\left(\frac{1}{n (\log n)^3}\right).$$

How many ways to choose such areas? At most n choices for starting vertex x .

Other vertices are in rectangle of area $(\log n) \Rightarrow$ at most $(\log n)^3$ choices for other vertices.

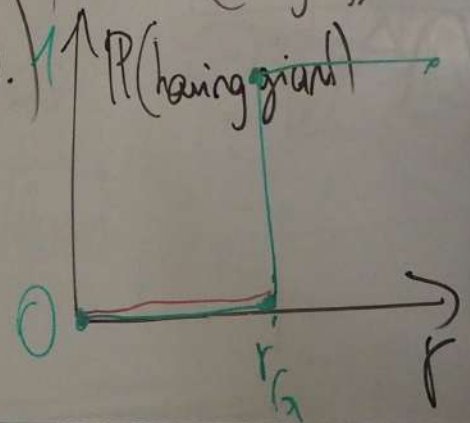
$$P(\exists \text{ such area empty}) = o(1)$$

Q: When is there already a giant component in EG_n in $[0, n^2]$ that is, a component of size $\Theta(n)$?

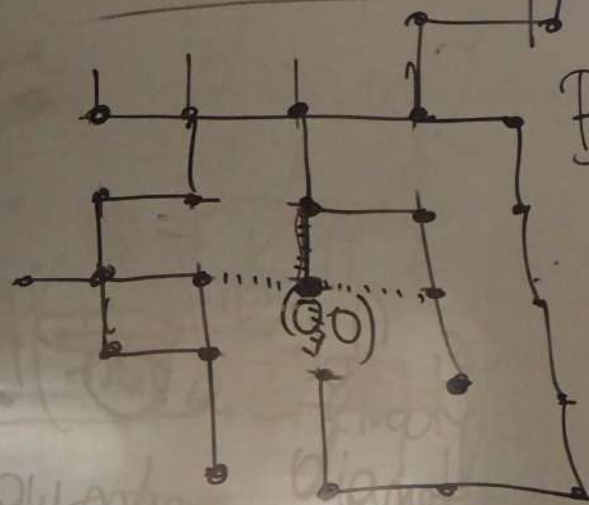
It will turn out that $C_1 \leq r_G \leq C_2$ for some constants C_1, C_2 .
(If $r \leq C_1$, then all components have whp at most $O((\log n)^2)$ vertices, and if $r \geq C_2$, then $\exists$ a component of size $\Theta(n)$.)

Open problem: find value of r_G

Here: show $C_1 \leq r_G \leq C_2$



Excursion to percolation on $\mathbb{Z}^2$:



Bond percolation: keep each edge independently with probability p , $0 \leq p \leq 1$

Big question: Is there an infinite cluster?

Is there positive probability that $(0,0)$ is in

an infinite component?

Define $\theta(p) = P_p[0 \text{ is in an infinite component}]$

$P_c = \inf \{p: \theta(p) > 0\}$

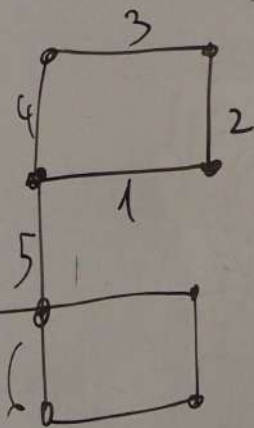
For $p < p_c$: $P(x \text{ is in an infinite component}) = 0 \quad \forall x$

$\Rightarrow$ by union bound, $P(\exists x \text{ in an infinite comp}) = 0$.

Theorem: $0 < p_c < 1$.

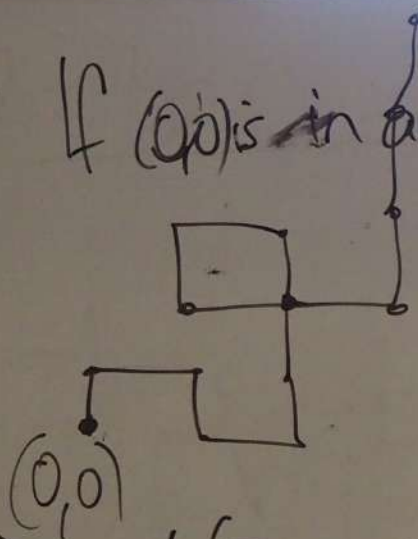
Proof: A self-avoiding path of length n :

path not reusing edges,
but can reuse vertices



Define $\mathcal{Q}_n$ to be the set of self-avoiding paths of length n starting at $(0,0)$

Obs: If $(0,0)$ is in an infinite component, then there exists a self-avoiding path of length n , $n \geq 0$ starting from $(0,0)$

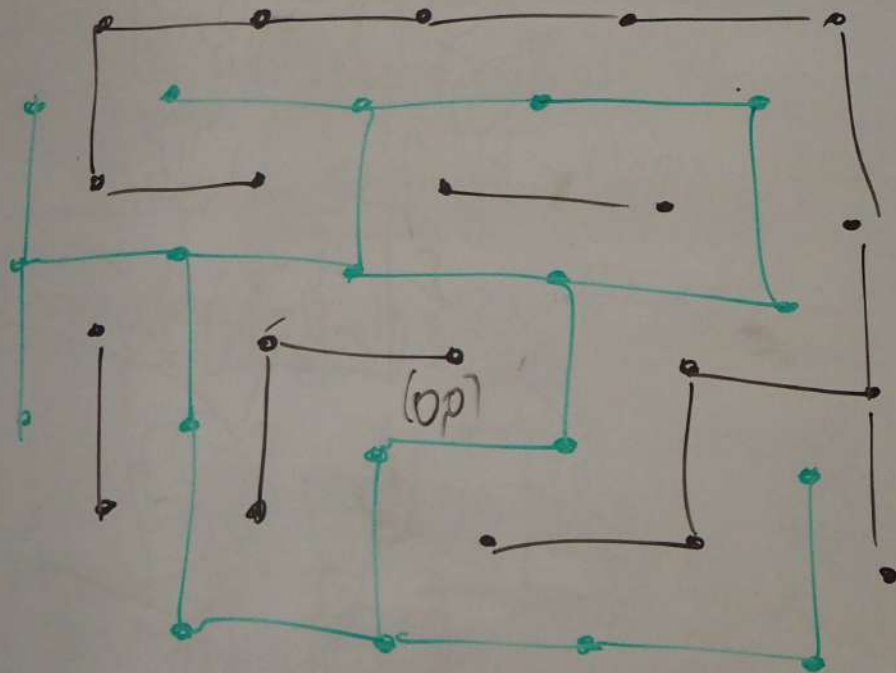


$$\theta(p) \leq P(\exists \text{ a self-avoiding path of length } n \text{ starting at } (0,0))$$

$$\leq \sum_{\gamma \in \mathcal{Q}_n} P(\text{each edge of the path of } \gamma \text{ is still present}) \leq |\mathcal{Q}_n| p^n \leq (4p)^n$$

If $p < \frac{1}{4}$, then this tends to 0 $\Rightarrow p_c \geq \frac{1}{4}$.

For an upper bound: we introduce the dual graph $(\mathbb{Z}^2)^*$.



$\mathbb{Z}^2$

$(\mathbb{Z}^2)^*$: dual edge is present
 $\Leftrightarrow$ original associated
 edge is not present

Obs: If $(0,0)$ is not in an infinite component, then, there must be a circuit in the dual graph surrounding $(0,0)$.

This circuit must intersect the point $(n+\frac{1}{2}, 0)$ for some $n \in \mathbb{N}$

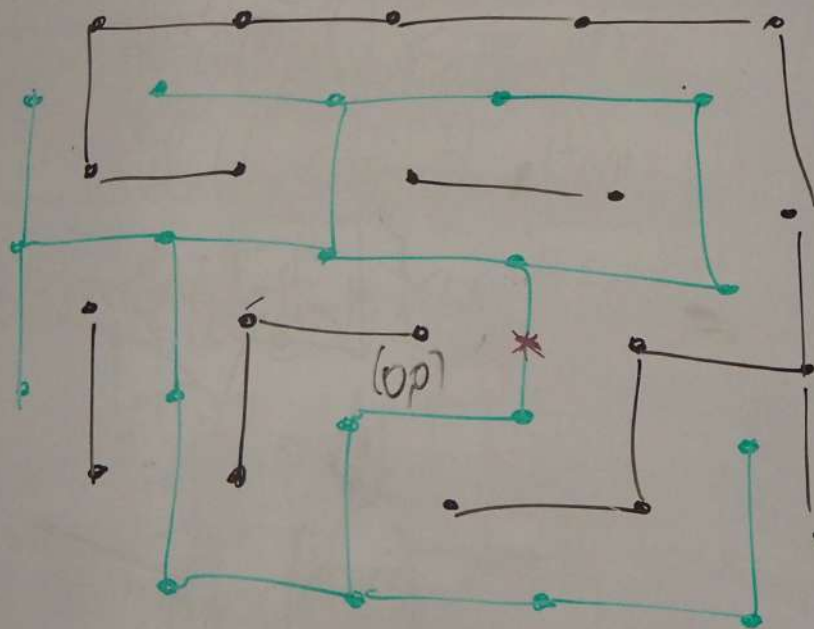
$$1 - \theta(\phi) \leq \sum_{n \geq 0} \mathbb{P}_p \left[\exists \text{ circuit in } (\mathbb{Z}^2)^* \text{ intersecting } (n+\frac{1}{2}, 0), \text{ surrounding the origin } (0,0) \right]$$

$$< \sum_{n \geq 0} \mathbb{P}_p \left[\exists \text{ path in } (\mathbb{Z}^2)^* \text{ of length at least } 2n+1 \text{ surrounding the origin } (0,0), \text{ intersecting } (n+\frac{1}{2}, 0) \right]$$

$$\leq \sum_{n \geq 0} (4(1-p))^{2n+4}$$

for p close to 1 (for ex. $p = 7/8$ still works)

this sum is strictly less than 1.



$$\Rightarrow 1 - \theta(p) < 1$$

$$\Rightarrow \theta(p) > 0$$

$$\Rightarrow p_c < 1$$

Remark: In fact $p_c = 1/2$