

Optimal Transport & Applications to Data Science

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Jan 17th, 2018

This talk

- Increased interest in ML community to use optimal transport (OT) methods.
- NIPS Tutorial on OT.
- New Book: Computational Optimal Transport (Peyre & Cuturi), FREE online.
- Classic Book: Cédric Villani.
- Hundreds of research papers in the past 5 years.
- Good opportunity to create interactions between mathematicians/theorists and data scientists/ML researchers.

Current (past) statistical paradigm

Maximum Likelihood estimator (popularized by Fisher)

$$\{x_i\}_{i=1}^N \sim p_{\bar{\theta}}(x), i.i.d., \theta_{MLE} = \arg \max_{\theta} \prod_{i=1}^N p_{\theta}(x_i)$$

- Nice interpretation.
- Nice properties: sufficiency, efficiency, etc
- Can be thought as minimization of a discrepancy. If p_{θ} is density of α_{θ} and $p_{\bar{\theta}}$ of $\hat{\beta}$ with empirical $\beta_n = \frac{1}{N} \sum_{i=1}^N \delta(x_i)$

$$\min_{\theta} \mathcal{L}_{MLE}(\alpha_{\theta}, \beta_n) \rightarrow KL(\alpha_{\theta}, \beta)$$

$$\mathcal{L}_{MLE}(\alpha_{\theta}, \beta_n) \equiv - \sum_{i=1}^N \log p_{\theta}(x_i) \rightarrow KL(\alpha_{\theta}, \beta)$$

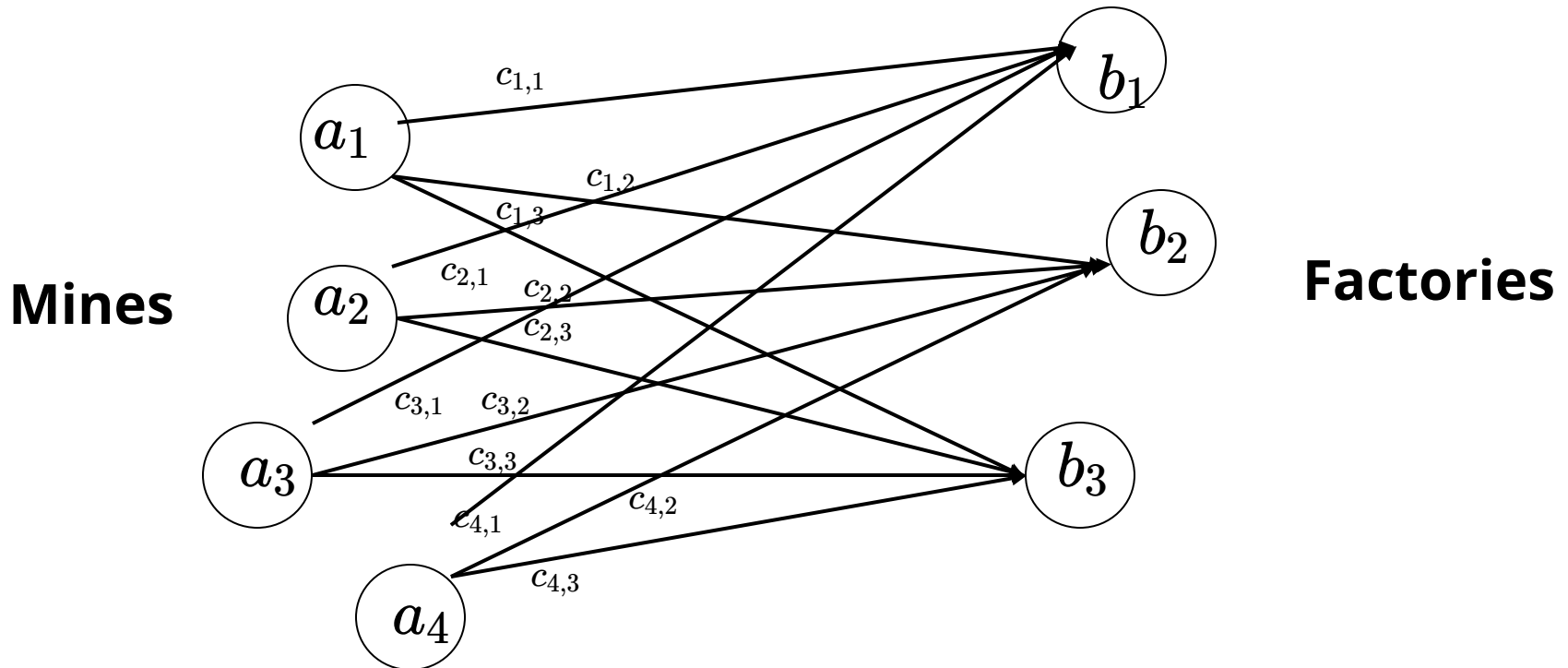
- It induces geometry, e.g. Fisher information metric, Amari 1985.
- However, it is not 'perfect'.

Optimal Transport

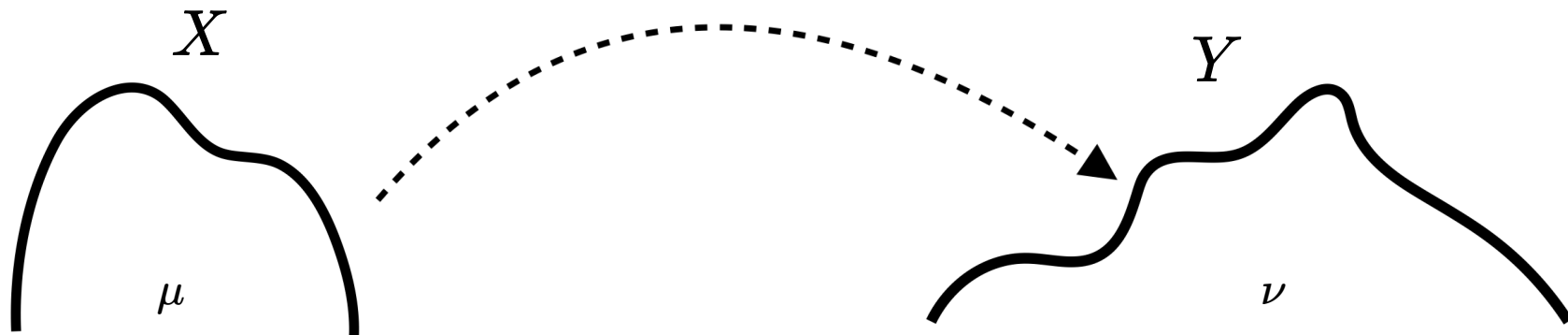
How to minimize the cost of transporting
all units in mines to factories?

$$\min p_{i,j} c_{i,j}$$

$$p_{i,j} \geq 0 \quad \sum_i p_{i,j} = a_i, \quad \sum_j p_{i,j} = b_j, \quad \sum_i a_i = \sum_j b_j$$



Optimal Transport



Monge Formulation

$$\inf_T \int_X c(x, T(x)) d\mu(x)$$

$$\text{s.t. } \nu = T\#\mu \text{ i.e. } \nu[B] = \mu(T^{-1}(B))$$

Kantorovich Formulation

$$\inf_{\pi} \int_{X \times Y} c(x, y) d\pi(x, y)$$

$$\text{s.t. } \pi \in \Pi[\mu, \nu], \text{ i.e. } \pi[X \times B] = \nu(B), \pi[A \times Y] = \mu(A)$$

Kantorovich Duality

Let X, Y be Polish spaces and $c : X \times Y \rightarrow \mathbb{R}_+ \cup \{+\infty\}$ be a lower-semi continuous cost function.

Define

$$I[\pi] = \int_{X \times Y} c(x, y) d\pi(x, y), \quad J(\phi, \psi) = \int_X \phi d\mu + \int_Y \psi d\nu$$

And

$$\Phi_c = \{(\phi, \psi) \in L^1(d\mu) \times L^1(d\nu) : \phi(x) + \psi(y) \leq c(x, y)\}$$

Then,

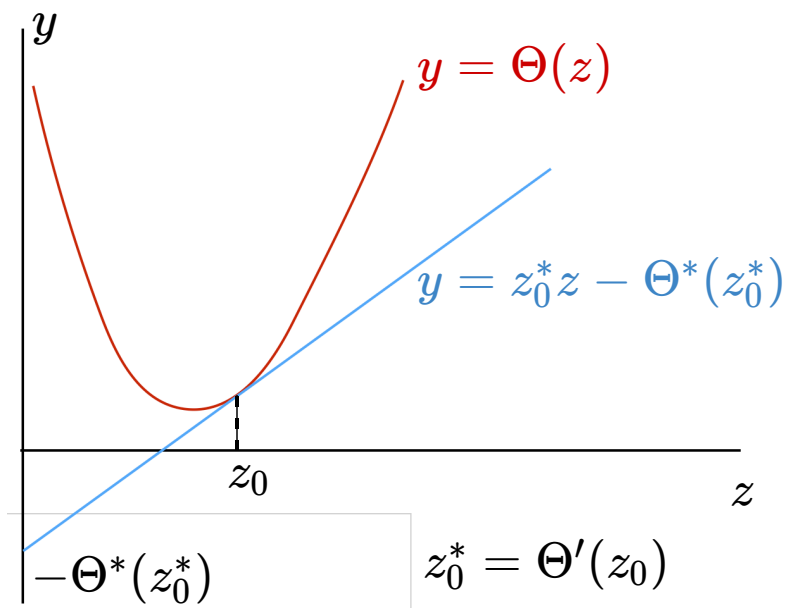
$$\inf_{\Pi(\mu, \nu)} I[\pi] = \sup_{\Phi_c} J(\phi, \psi)$$

Legendre-Fenchel transform

Let $\Theta(\cdot)$ be convex, on a n.v.s E with values in $\mathbb{R} \cup \{+\infty\}$

For values in the topological dual E^* define the L-F conjugate as:

$$\Theta^*(z^*) = \sup_{z \in E} [\langle z^*, z \rangle - \Theta(z)]$$

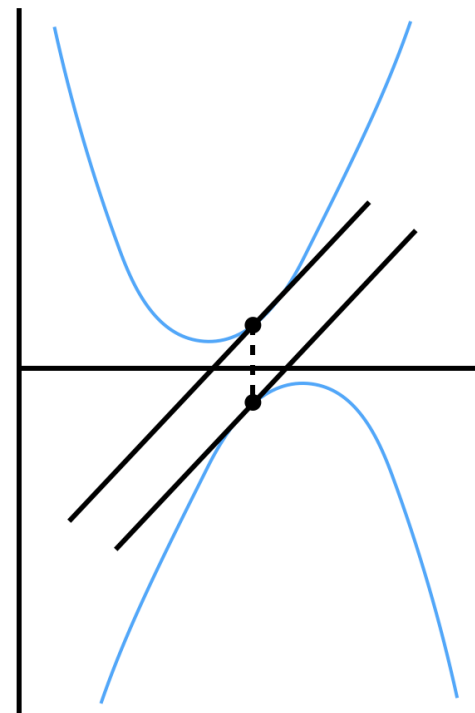


Fenchel-Rockafellar theorem

Let $\Theta(\cdot), \Xi(\cdot)$ be convex. If for some $z_0 \in E$ $\Theta(\cdot)$ is continuous at z_0 and $\Theta(z_0) < +\infty, \Xi(z_0) < +\infty$

Then,

$$\inf_{z \in E} [\Theta + \Xi] = \max_{z^* \in E^*} [-\Theta^*(-z^*) - \Xi^*(z^*)]$$



Proof sketch

1. Start with X, Y compact, c continuous and use $E = C_b(X \times Y)$ (here $E^* = M(X \times Y)$ regular Bore measures). Show it holds there by Fenchel-Rockafellar theorem.
2. $\Pi(\mu, \nu)$ is compact; infimum is attained.
3. Relax compactness using tightness of the minimizer π^* and a truncation argument. Notice one can use $E = L^1(d\mu) \times L^1(d\nu)$.

Kantorovich- Rubinstein Theorem

Let $X = Y$ be a Polish space and $c(x, y) = d(x, y)$ a lower semi continuous metric.

Define

$$\mathcal{T}_d = \inf_{\pi \in \Pi(\mu, \nu)} \int_{X \times X} d(x, y) d\pi(x, y)$$

And

$$\|\phi\|_{Lip} \equiv \sup_{x \neq y} \frac{|\phi(x) - \phi(y)|}{d(x, y)}$$

Then,

$$\mathcal{T}_d = \sup \left\{ \int_X \phi d(\mu - \nu); \phi \in \cap L^1(d|\mu - \nu|); \|\phi\|_{Lip} \leq 1 \right\}$$

Relevant case: $d(x, y) = \|x - y\|$.

Proof sketch

1. For ϕ bounded define the pair of c-concave conjugate pair (ϕ^c, ϕ^{cc}) as:

$$\phi^c(y) = \inf_{x \in X} [c(x, y) - \phi(x)], \phi^{cc}(y) = \inf_{x \in Y} [c(x, y) - \phi^c(x)]$$

2. Notice one can replace (ϕ, ψ) by (ϕ^{cc}, ϕ^c) and $\phi^c = -\phi^{cc}$

3.
$$\sup_{\phi \in L^1(d\nu)} J(-\phi^d, \psi^d) = \sup_{\|\phi\|_{Lip} \leq 1} J(\phi, -\phi)$$

The case $c(x, y) = \|x - y\|^2$

Knott-Smith optimality criterion

$\pi \in \Pi(\mu, \nu)$ is optimal if there exists a convex lower semi-continuous $\phi(x)$ satisfying,

For $d\pi$ almost all (x, y) , $y \in \partial\phi(x)$

Moreover, (ϕ, ϕ^*) minimizes

$$\inf\left\{\int_{\mathbb{R}^n} \phi d\mu + \int_{\mathbb{R}^n} \psi d\nu; \forall (x, y), x \cdot y \leq \phi(x) + \psi(y)\right\}$$

Brenier's theorem

If μ does not give mass to Lebesgue-negligible sets then the unique optimal π satisfies.

$$d\pi(x, y) = d\mu(x)\delta[y = \nabla\phi(x)]$$

$\nabla\phi(x)$ is the unique gradient of a convex function s.t $\nu = \nabla\phi\#\mu$

The metric side of OT

Let X be a Polish space endowed with a distance d and $p \geq 0$

Define

$$\mathcal{T}_p = \inf_{\pi \in \Pi(\mu, \nu)} \int_{X \times X} d^p(x, y) d\pi(x, y)$$

Then, if $p \in [1, \infty)$, $W_p = \mathcal{T}_p^{1/p}$ defines a metric on $\mathcal{P}_p(X)$

(p-moment bounded)

If $p \in [0, 1)$, $W_p = \mathcal{T}_p$ defines a metric on $\mathcal{P}_p(X)$

Wasserstein distances are distances between distributions and metrize weak convergence.

Displacement interpolation

Time dependent (Monge) optimal transport

$$\inf\left\{\int_X C[(T_t x)]_{0 \leq t \leq 1} d\mu(x); T_0 = Id, T_1 \# \mu = \nu\right\}$$

C is consistent with original Monge if optimal T_t gives rise to optimal $T = T_1$. e.g., $c(x, y) = \inf\{C[(z_t)_{0 \leq t \leq 1}], z_0 = x, z_1 = y\}$

$$C[(z_t)_{0 \leq t \leq 1}] = \int_0^1 |\dot{z}_t|^2 dt \rightarrow c(x, y) = |x - y|^2$$

In this case (**McCann displacement interpolation**)

$$T_t(x) = (1 - t)x + t\nabla\phi(x) \leftrightarrow \rho_t = [(1 - t)Id + t\nabla\phi] \# \mu$$

$$W_2(\mu, \rho_t) = tW_2(\mu, \nu)$$

Constant velocity geodesic.

Computational Optimal transport

OT problem states: $d_c(\mu, \nu) = \inf_{\pi \in \Pi(\mu, \nu)} E_\pi(c(x, y))$

In practice, transportation between discrete measures

$$\mu = \sum_{i=1}^n u_i \delta_{x_i}, \nu = \sum_{j=1}^m v_j \delta_{y_j}$$

The problem becomes

$$d_c(\mu, \nu) = \inf_{P \in U(u, v)} \langle P, C \rangle$$

$$\langle P, C \rangle = \sum_{i,j} P_{i,j} C_{i,j}, \quad C_{i,j} = c(x_i, y_j)$$

In the transportation polytope

$$U(u, v) = \{P \in \mathcal{R}_+^{n \times m}; P1_m = u, P^T 1_n = v\}$$

Entropy regularized Optimal Transport. Why?

The discrete OT is solved by combinatorial optimization algorithm with $O(n^3)$ Perhaps better approximate solutions?
 $O(n^3 \log(n))$

The operator $C \mapsto \arg \min_{\gamma} \langle C, \gamma \rangle$ is not differentiable. Too bad for automatic differentiation libraries.

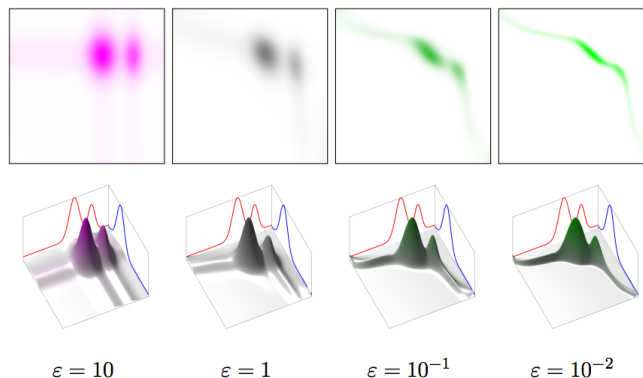
Entropy regularized Optimal Transport

Add entropy term $h(P) = -\sum_{i,j} P_{i,j} \log P_{i,j}$ Cuturi, 2012

$$d_c^\lambda(\mu, \nu) = \inf_{P \in U(u,v)} \langle P, C \rangle - \lambda h(P) \quad \text{Cuturi function}$$

$$= \inf_{P \in U^\alpha(u,v)} \langle P, C \rangle, \quad U^\alpha(u,v) = \{P \in U(u,v) : KL(P \| uv^\top) \leq \alpha\}$$

i.e., enforce contingency tables to be close to the independent



From Computational OT book

Sinkhorn algorithm

First order optimality condition

$$P_{i,j}^\lambda = \exp(-1/2 - \alpha_i/\lambda) \exp(-c_{i,j}/\lambda) \exp(-1/2 - \beta_j/\lambda)$$

with α, β Lagrange multipliers i.e.,

$$P^\lambda = \text{diag}(a) K^\lambda \text{diag}(b), \quad a, b > 0, K^\lambda = \exp(-C/\lambda)$$

By Sinkhorn's theorem, there is a unique $P^\lambda \in U(u, v)$ as above and can be obtained by Sinkhorn fixed point iterations (matrix scaling)

$$a^{l+1} = \frac{u}{K b^l}, \quad b^{l+1} = \frac{v}{K^\top a^{l+1}}$$

Other relations

Bregman projections

$$\mathcal{C}_u = \{P \in \mathcal{R}_+^{n \times m}; P1_m = u\}, \mathcal{C}_v = \{P \in \mathcal{R}_+^{n \times m}; P^\top 1_n = v\}$$

$$d_c^\lambda(\mu, \nu) = \inf_{P \in U(u, v)} KL(P \| K^\lambda)$$

$$P^{l+1} = Proj_{\mathcal{C}_u}^{KL}(P^l), P^{l+2} = Proj_{\mathcal{C}_v}^{KL}(P^{l+1})$$

Duality

$$d_c^\lambda(\mu, \nu) = \max_{f \in \mathbb{R}^n, g \in \mathbb{R}^m} \langle f, u \rangle + \langle g, v \rangle - \lambda \langle e^{f/\lambda}, K e^{g/\lambda} \rangle$$

Convergence to actual solution. Cominetti & San Martin 1994

Reverse mode automatic differentiation

Solution can be written as $\langle \text{diag}(a^l) K^\lambda \text{diag}(b^l), C \rangle$

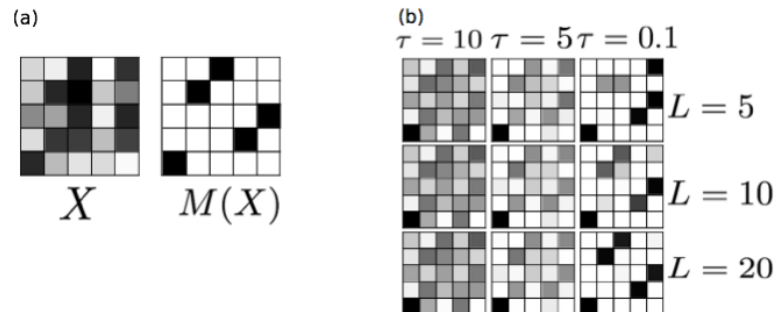
Differentiate this expression with respect to relevant parameters

The case of permutations

- If $m = n$ one solves a linear program on the Birkhoff polytope. $\mathcal{B}$
- Sinkhorn algorithm becomes iterative row and column scaling of a matrix, Sinkhorn operator $S(\cdot)$.
- Entropy regularization allow us to conceive soft matchings

$$S(C/\lambda) = \arg \max_{P \in \mathcal{B}} \langle P, C \rangle + \lambda h(P)$$

- In the limit become hard permutations
 $M(C) \equiv \arg \max_{P \in \mathcal{B}} \langle P, C \rangle = \lim_{\lambda \rightarrow 0^+} S(C/\lambda)$
- Analogy with softmax



Complexity of Sinkhorn algorithm

Active area of research. Sinkhorn algorithm works in near-linear time

Theorem (Altschuler, Weed, Rigollet, 2017). Sinkhorn algorithm gives a solution $\hat{P}$ satisfying

$$\langle \hat{P}, C \rangle \leq \min_{P \in U(u,v)} \langle P, C \rangle + \epsilon$$

in $O(n^2 L^4 \log(n)^2 \epsilon^{-3})$ operations.

Here, $\epsilon = \lambda \log n$, $\epsilon' = \epsilon / (4 \|C\|_\infty)$ $\|C\|_\infty \leq L$

Sinkhorn iterations are defined with a tolerance ϵ'

Wasserstein Barycenters

$$\bar{\mu} = \arg \min_u \sum_{i=1}^N \lambda_i W_2^2(\mu, \nu_i)$$

Generalizes McCann displacement formula (Agueh and Carlier, 2010)

Relates to Multi-Marginal problem (Agueh and Carlier, 2010)

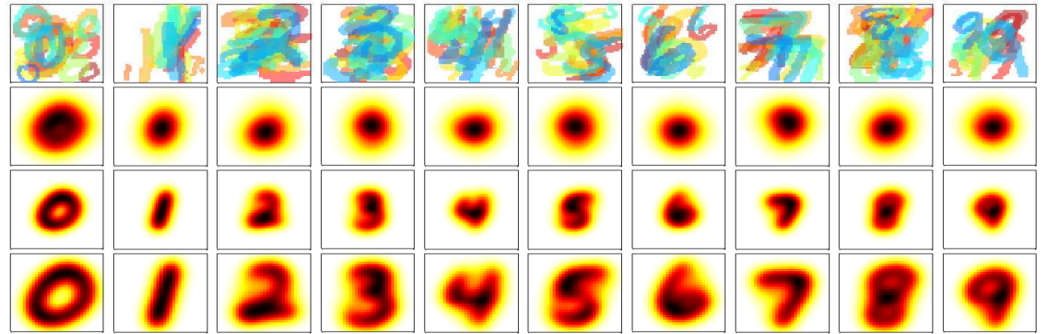
$$\inf_{\pi \in \Pi(\mu_1, \dots, \mu_p)} \int_{\mathbb{R}^d} \left(\sum_{i=1}^p \frac{\lambda_i}{2} |x_i - T(x)|^2 \right) d\pi(x_1, \dots, x_p)$$

$$\bar{\mu} = \left(\sum_{i=1}^p T_i^j \right) \# \nu_j$$

Wasserstein barycenters

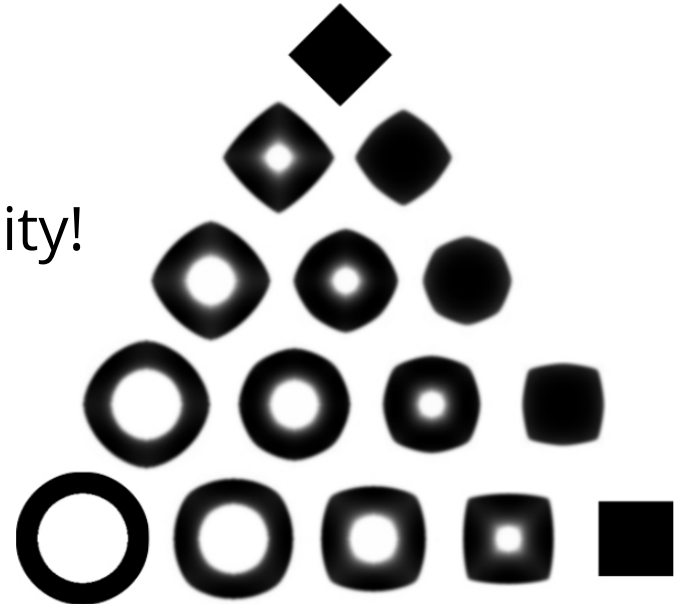
Cuturi and Doucet, 2014

Images as histograms



Benamou et al 2014

Use entropic regularization for scalability!



Wasserstein Principal Geodesic Analysis

Fletcher et al, 2004. Extension of PCA to the manifold setting

Algorithm 2: Principal Geodesic Analysis

Input: $x_1, \dots, x_N \in M$

Output: Principal directions, $v_k \in T_\mu M$

Variances, $\lambda_k \in \mathbb{R}$

$\mu =$ intrinsic mean of $\{x_i\}$ (Algorithm 1)

$u_i = \text{Log}_\mu(x_i)$

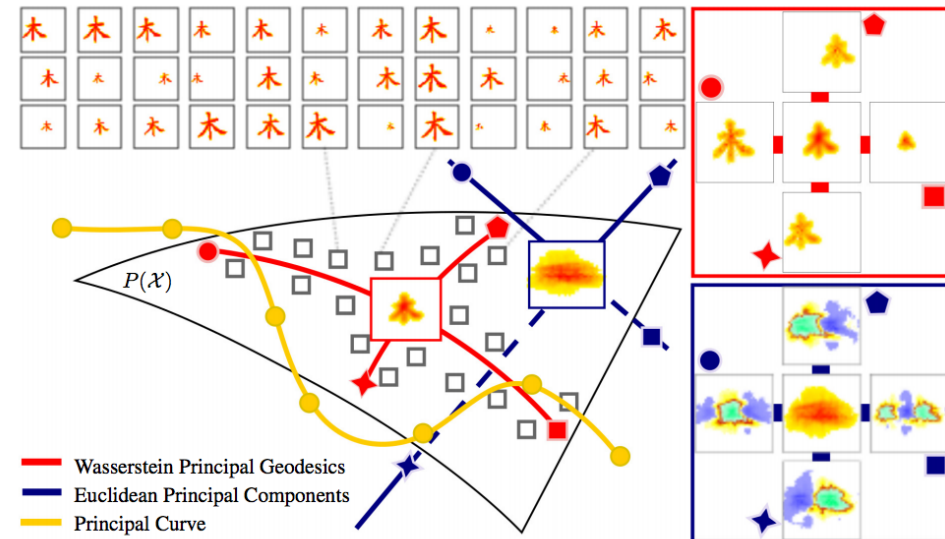
$S = (1/N) \sum_{i=1}^N u_i u_i^T$

$\{v_k, \lambda_k\} =$ eigenvectors/eigenvalues of S .

Wassersteing case

Bigot, Gouet, Klein, Lopez 2015.

Scaled by Seguy and Cuturi, 2015



Generative models

How to estimate θ in a model defined as follows: Sample from a reference (noise) distribution $z \sim \zeta \in \mathcal{P}(\mathcal{Z})$, then $x = h_\theta(z) \in X$

$h_\theta : \mathcal{Z} \rightarrow X$ is usually a complex nonlinear function between two high-dimensional spaces. $\alpha_\theta = h_{\theta\#}\zeta$ doesn't have a tractable density and cannot use MLE.

Instead, consider the problem (β is the true distribution of data.)

$$\begin{aligned} \min_{\theta} W_1(\alpha_\theta, \beta) &= \min_{\theta} \inf_{P \in \Pi(\alpha_\theta, \beta)} E(\|x - y\|) \\ &= \min_{\theta} \sup_{f \in Lip} E_{\alpha_\theta}(f(x)) - E_{\beta}(f(x)) \end{aligned}$$

Solve minimax problem using $\nabla_{\theta} W_1(\alpha_\theta, \beta) = E_{z \sim \zeta}(\nabla_{\theta} f(h_{\theta}(z)))$

Wasserstein GAN, Arjovsky et al 2017

More generative models

Alternative formulation

Genevay et al 2017 uses primal formulation

$$\min_{\theta} W_1(\alpha_{\theta}, \beta) = \min_{\theta} \inf_{P \in \Pi(\alpha_{\theta}, \beta)} E(\|x - y\|)$$

Also, use entropy regularization to enable automatic differentiation.

Alternative distances, Integral probability metrics

$$\mathcal{E}(\alpha, \beta)^2 \equiv 2E_{x \sim \alpha, y \sim \beta}(\|x - y\|) - E_{x \sim \alpha, x' \sim \alpha}(\|x - x'\|) - E_{y \sim \beta, y' \sim \beta}(\|y - y'\|)$$

$$\mathcal{E}(\alpha, \beta)^2 = \sup_{f \in \mathcal{H}, \|f\|_{\mathcal{H}} \leq 1} E_{\alpha}(f(x)) - E_{\beta}(f(x))$$

$$D_{TV}(\alpha, \beta) = \sup_{A \in \mathcal{F}} |\alpha(A) - \beta(A)| = \sup_{f \in C(X, [0,1])} E_{\alpha}(f(x)) - E_{\beta}(f(x))$$

More generative models

Sample complexity

Computations involve samples from the model and data.

$$E(|W_1(\alpha_n, \beta_n) - W_1(\alpha, \beta)|) = O(n^{-1/d})$$

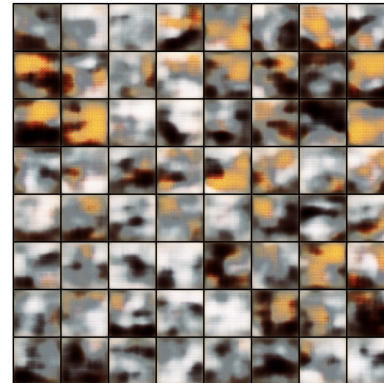
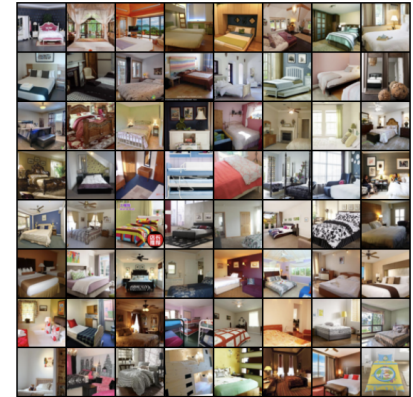
d = data dimension

$$E(|\mathcal{E}_1^2(\alpha_n, \beta_n) - \mathcal{E}^2(\alpha, \beta)|) = O(n^{-1})$$

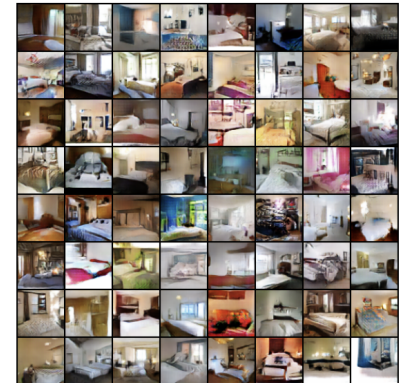
This would be catastrophic, but in practice it is not.

Why? study differences in induced geodesic structure

Bottou et al, 2017



Generated by the ED trained model



Generated by the WD trained model

Robustness to adversarial attacks

Goal: quantify robustness to adversarial examples and design robust algorithms

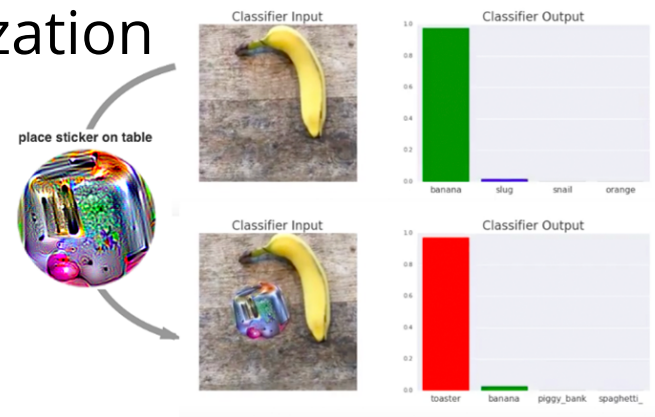
Classical approach: Empirical Risk Minimization

$$\min_{\theta} \frac{1}{N} \sum_{i=1}^N l(\theta, x_i) = \min_{\theta} E_{P_n}(l(\theta, x))$$

'Modern' approach: distributionally robust optimization

$$\min_{\theta} \sup_{P: W(P, P_n) \leq \rho} E_P(l(\theta, x))$$

Sinha, Namkook, Duchi 2017. Robust algorithm, and certificate
For the true P_0



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